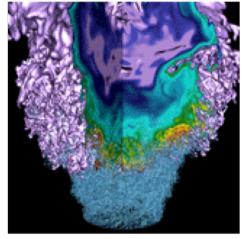
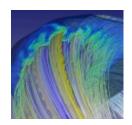
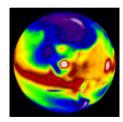
Scalable HDF5

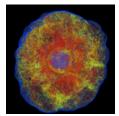


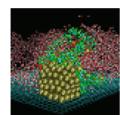












Quincey Koziol

ATPESC August 3, 2018 koziol@lbl.gov





Why HDF5?

Have you ever asked yourself:

- How will I deal with one-file-per-processor in the petascale era?
- Do I need to be an "MPI and Lustre pro" to do my research?
- Where is my checkpoint file?

HDF5 hides all complexity so you can concentrate on Science

Optimized I/O to single shared file*

^{*} Explorations for "multi-file" HDF5 storage are underway as well.





Goal

Introduce you to HDF5

- HDF5 Overview
- HDF5 Programming Overview
- Intro to Scalable HDF5





WHAT IS HDF5?





What is HDF5?

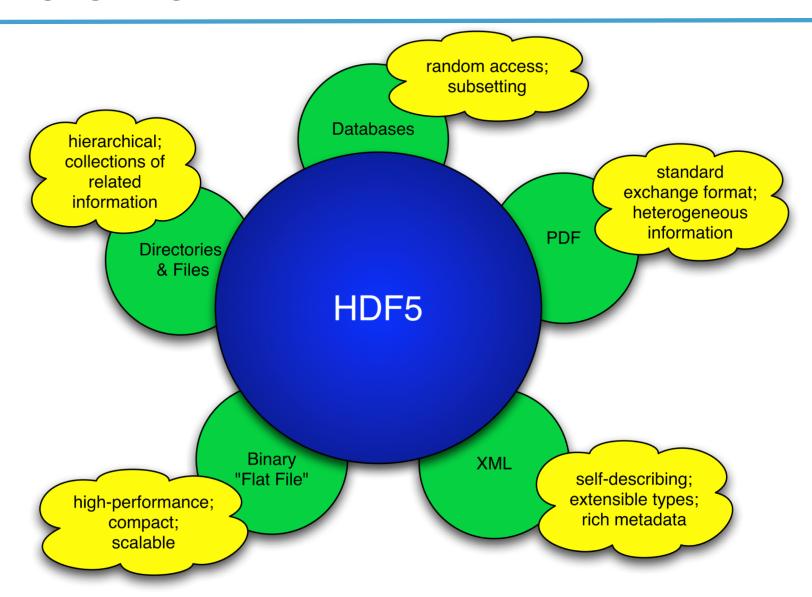
HDF5 == Hierarchical Data Format, v5

- Open file format
 - Designed for high volume or complex data
- Open source software
 - Works with data in the format
- An extensible data model
 - Structures for data organization and specification





HDF5 is like ...





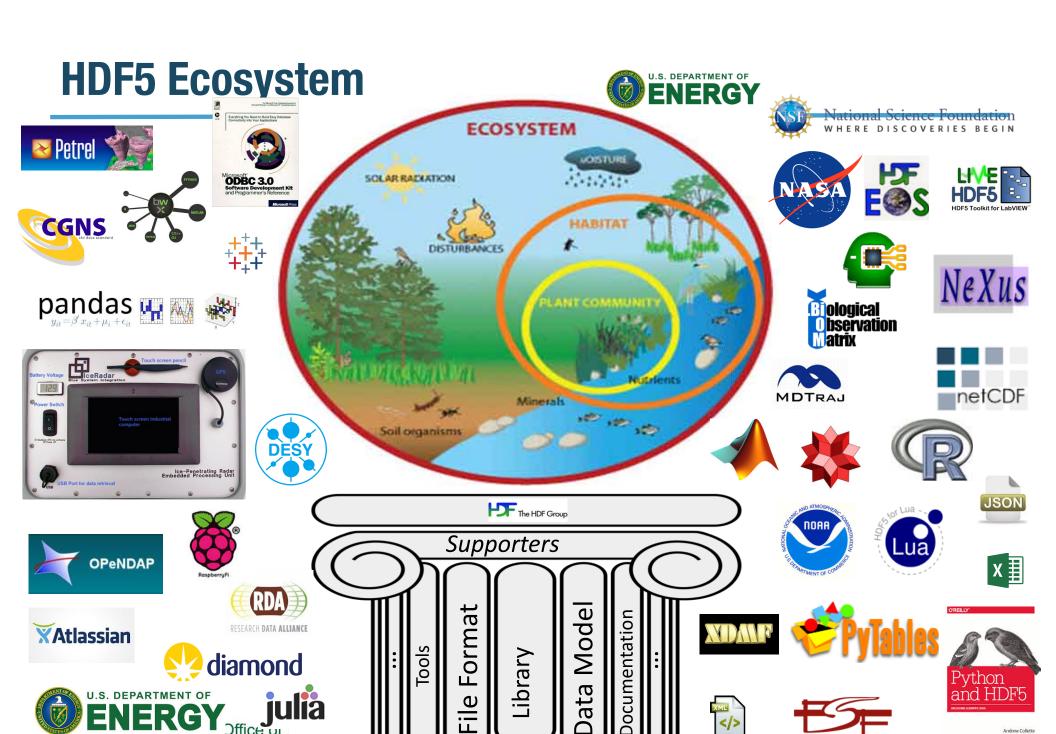


HDF5 is designed ...

- for high volume and/or complex data
- for every size and type of system (portable)
- for flexible, efficient storage and I/O
- to enable applications to evolve in their use of HDF5 and to accommodate new models
- to support long-term data preservation







Science

BERKELEY LAB
Lawrence Berkeley National Laboratory

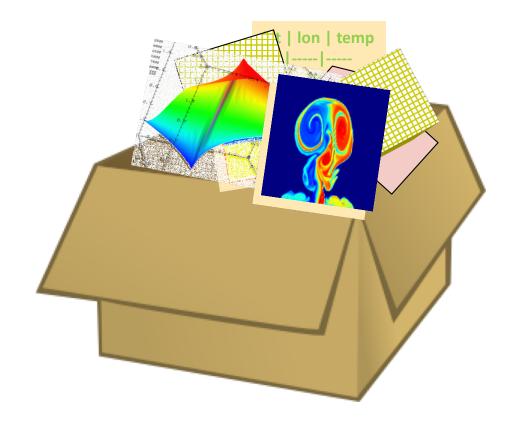
HDF5 DATA MODEL





HDF5 File

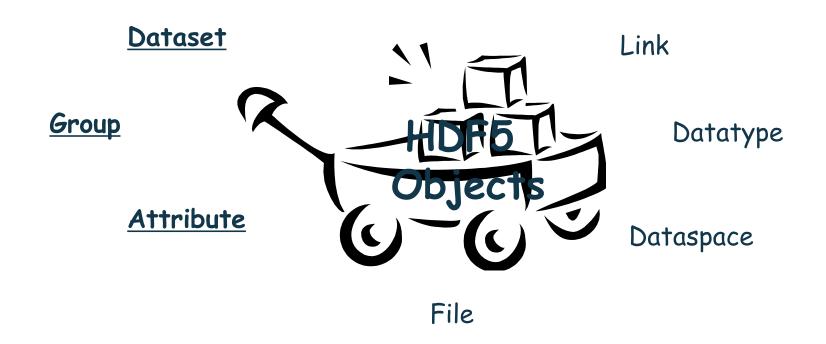
An HDF5 file is a **container** that holds data objects.







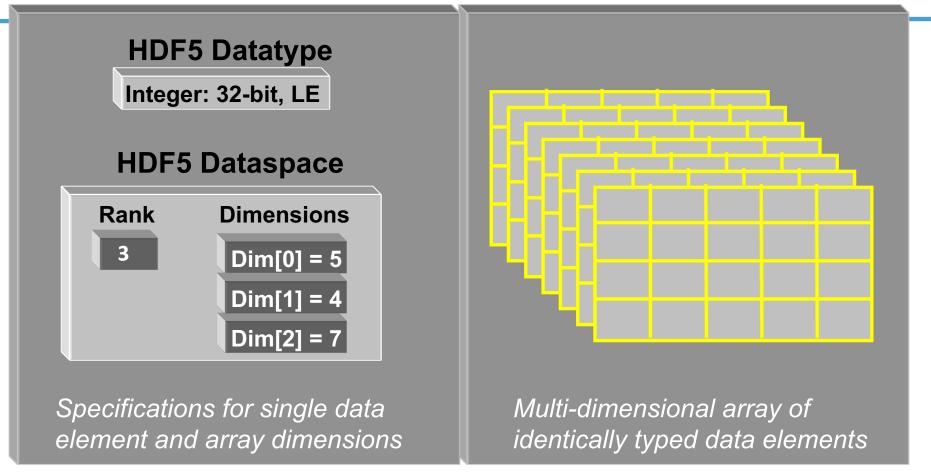
HDF5 Data Model







HDF5 Dataset



- HDF5 datasets organize and contain data elements.
 - HDF5 datatype describes individual data elements.
 - HDF5 dataspace describes the logical layout of the data elements.





HDF5 Dataspace

- Describes the logical layout of the elements in an HDF5 dataset
 - NULL
 - no elements
 - Scalar
 - single element
 - Simple array (most common)
 - multiple elements organized in a rectangular array
 - rank = number of dimensions
 - dimension sizes = number of elements in each dimension
 - maximum number of elements in each dimension
 - » may be fixed or unlimited



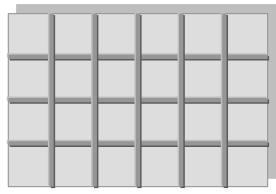


HDF5 Dataspace

Two roles:

<u>Spatial information for Datasets and Attributes</u>

- Rank and dimensions
- Permanent part of object definition



Rank = 2Dimensions = 4x6

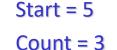
Partial I/O: Dataspace and selection describe application's data buffer and data elements participating in I/O



Rank = 1

Dimension = 10







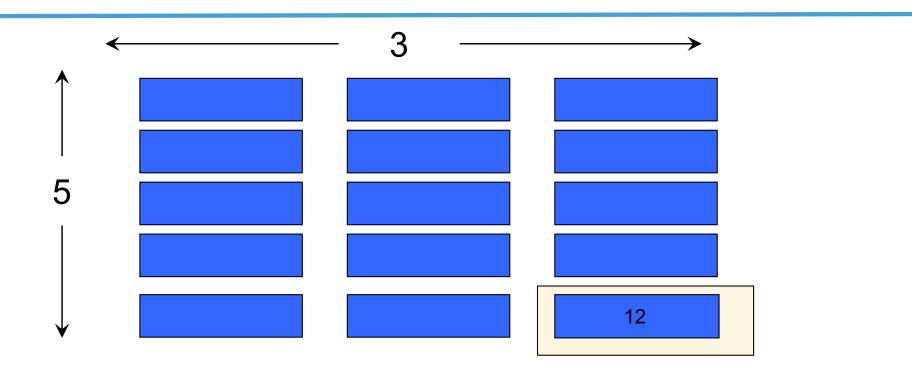
HDF5 Datatypes

- Describe individual data elements in an HDF5 dataset
- Wide range of datatypes supported
 - Integer
 - Float
 - Enum
 - Array (similar to matrix)
 - User-defined (e.g., 12-bit integer, 16-bit float)
 - Variable-length types (e.g., strings, vectors)
 - Compound (similar to C structs)
 - More ...





HDF5 Dataset



Datatype: 32-bit Integer

Dataspace: Rank = 2

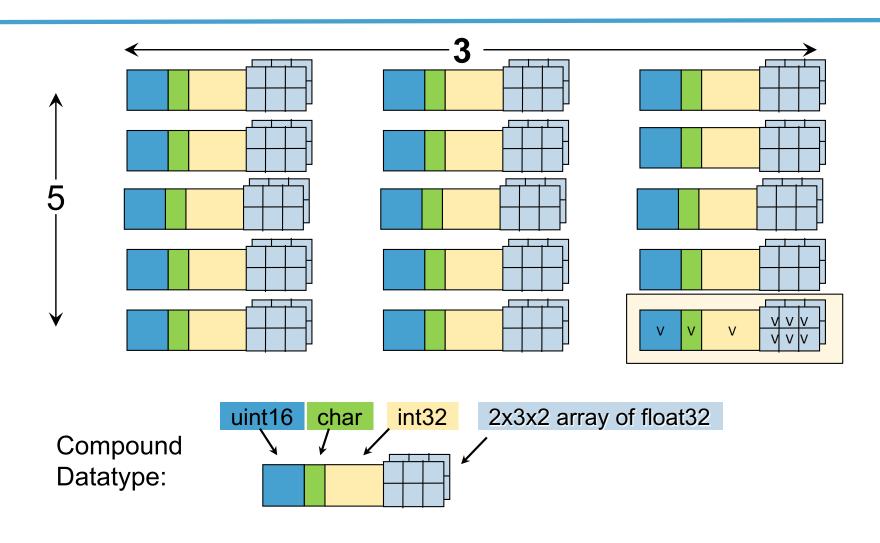
Dimensions = 3×5

Note that this is declared in C as: "array[5][3]" and as "array(3)(5)" in FORTRAN





HDF5 Dataset with Compound Datatype



Dataspace: Rank = 2

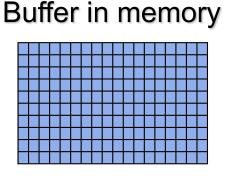
Dimensions = 5×3



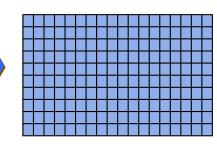


How are data elements stored? (1/2)

Contiguous (default)

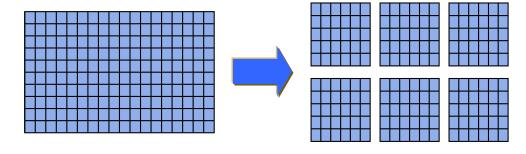


Data in the file



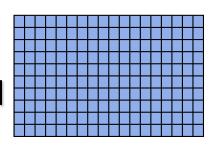
Data elements stored physically adjacent to each other

Chunked

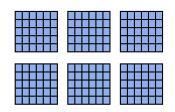


Better access time for subsets; extendible

Chunked & Compressed







Improves storage efficiency, transmission speed





How are data elements stored? (2/2)

Buffer in memory

Data in the file

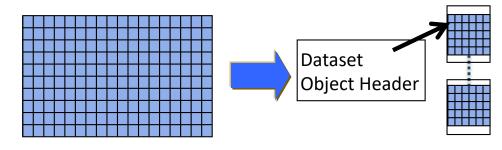
Compact





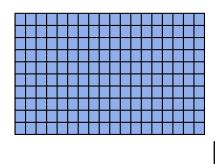
Data elements stored directly within object's metadata

External



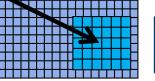
Data elements stored outside the HDF5 file, possibly in another file format

Virtual



Data elements actually stored in "source datasets", using selections to map







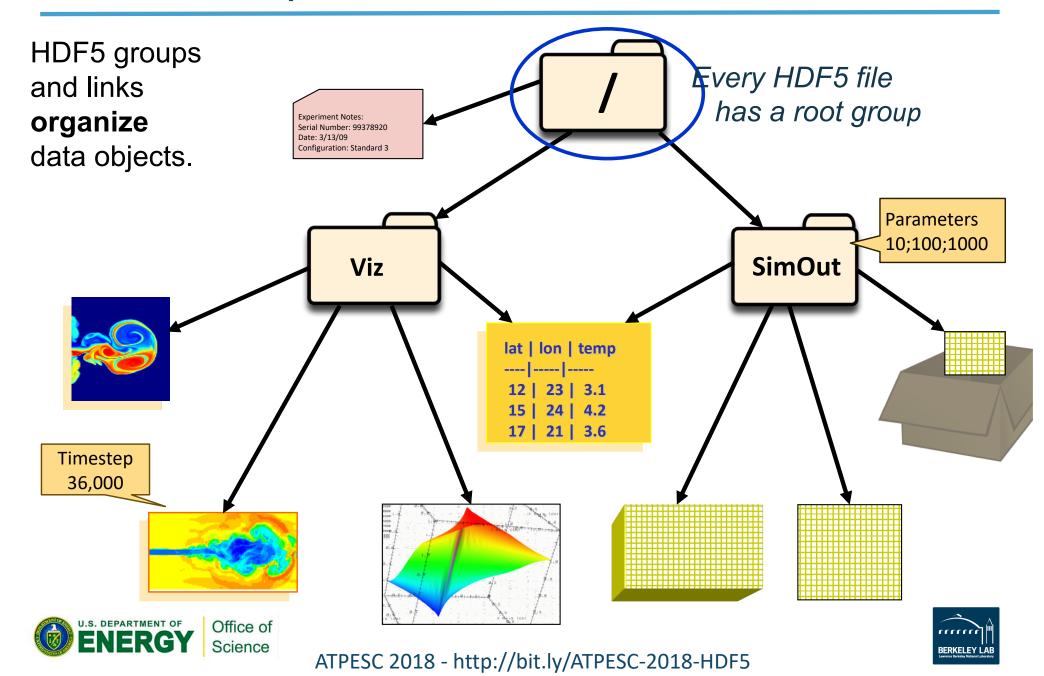
HDF5 Attributes

- Attributes "decorate" HDF5 objects
- Typically contain user metadata
- Similar to Key-Values:
 - Have a unique <u>name</u> (for that object) and a <u>value</u>
- Analogous to a dataset
 - "Value" is an array described by a datatype and a dataspace
 - Do not support partial I/O operations; nor can they be compressed or extended





HDF5 Groups and Links



HDF5 SOFTWARE





HDF5 Home Page

HDF5 home page: http://hdfgroup.org/HDF5/

Latest release: HDF5 1.10.2 (1.10.2 coming August 2018)

HDF5 source code:

- Written in C, and includes optional C++, Fortran APIs, and High Level
 APIs
- Contains command-line utilities (h5dump, h5repack, h5diff, ..) and compile scripts

HDF5 pre-built binaries:

- When possible, include C, C++, Fortran, and High Level libraries.
 Check ./lib/libhdf5.settings file.
- Built with and require the SZIP and ZLIB external libraries





Useful Tools For New Users

h5dump:

Tool to "dump" or display contents of HDF5 files

h5cc, h5c++, h5fc:

Scripts to compile applications

HDFView:

Java browser to view HDF5 files

http://www.hdfgroup.org/hdf-java-html/hdfview/

HDF5 Examples (C, Fortran, Java, Python, Matlab, ...) https://www.hdfgroup.org/HDF5/examples/

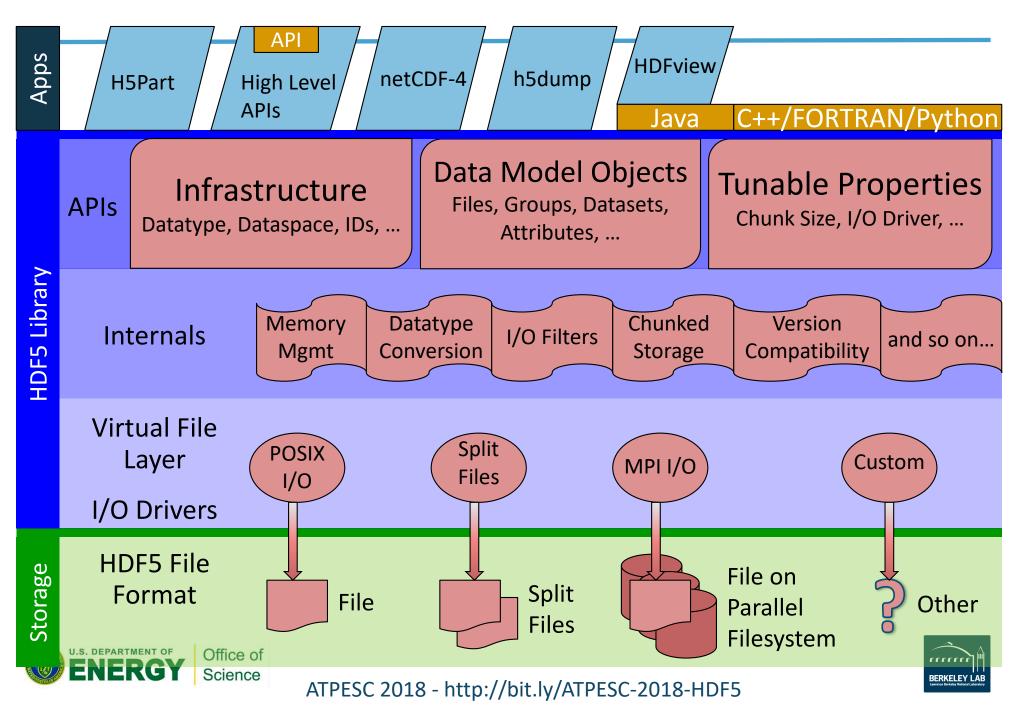


HDF5 PROGRAMMING MODEL AND API





HDF5 Software Layers & Storage



The General HDF5 API

- C, FORTRAN, Java, C++, and .NET bindings
- IDL, MATLAB, Python (H5Py, PyTables)
- C routines begin with prefix: H5?

? is a character corresponding to the type of object the function acts on

Example Functions:

H5D: Dataset interface e.g., **H5Dread**

H5F: File interface *e.g.,* **H5Fopen**

H5S: dataSpace interface e.g., H5Sclose





The HDF5 API

- For flexibility, the API is extensive
 - √ 300+ functions



- This can be daunting... but there is hope

- ✓ A few functions can do a lot
- √ Start simple
- ✓ Build up knowledge as more features are needed





General Programming Paradigm

- Object is opened or created
- Object is accessed, possibly many times
- Object is closed

- Properties of object are optionally defined
 - ✓ Creation properties (e.g., use chunking storage)
 - ✓ Access properties





Basic Functions

H5Fcreate (H5Fopen)

create (open) File

H5Screate_simple/H5Screate create dataSpace

H5Dcreate (H5Dopen) create (open) Dataset

H5Dread, H5Dwrite access Dataset

H5Dclose close Dataset

H5Sclose close dataSpace

H5Fclose close File





Other Common Functions

DataSpaces: H5Sselect_hyperslab (Partial I/O)

H5Sselect_elements (Partial I/O)

H5Dget_space

DataTypes: H5Tcreate, H5Tcommit, H5Tclose

H5Tequal, H5Tget_native_type

Groups: H5Gcreate, H5Gopen, H5Gclose

Attributes: H5Acreate, H5Aopen_name,

H5Aclose, H5Aread, H5Awrite

Property lists: H5Pcreate, H5Pclose

H5Pset_chunk, H5Pset_deflate





Tools

PARALLEL HDF5





Terminology

- DATA problem-size data, e.g., large arrays
- METADATA is an overloaded term
- In this presentation:
 - Metadata "=" HDF5 metadata
 - For each piece of application metadata, there are many associated pieces of HDF5 metadata
 - There are also other sources of HDF5 metadata





Why Parallel HDF5?

- Take advantage of high-performance parallel I/O while reducing complexity
 - Add a well-defined layer to the I/O stack
 - Keep the dream of a single or a few shared files alive
 - "Friends don't let friends use one file per process!"
- Make performance portable





What We'll Cover Here

- Parallel vs. serial HDF5
- Implementation layers
- HDF5 files (= composites of data & metadata) in a parallel file system
- PHDF5 I/O modes: collective vs. independent
- Data and metadata I/O





What We Won't Cover

- Consistency semantics
- Virtual Object Layer (VOL)
- Automatic tuning
- Single Writer / Multiple-Reader (SWMR)
- Virtual Datasets (VDS)
- Asynchronous I/O
- Independent Metadata Modification
- •

Come see me this evening or after the presentation!





(MPI-)Parallel vs. Serial HDF5

- PHDF5 allows multiple MPI processes in an MPI application to perform I/O to a single HDF5 file
- Uses a standard parallel I/O interface (MPI-IO)
- Portable to different platforms
- PHDF5 files <u>ARE</u> HDF5 files conforming to the <u>HDF5</u> file format specification
- The PHDF5 API consists of:
 - The standard HDF5 API
 - A few extra knobs and calls
 - A parallel "etiquette"





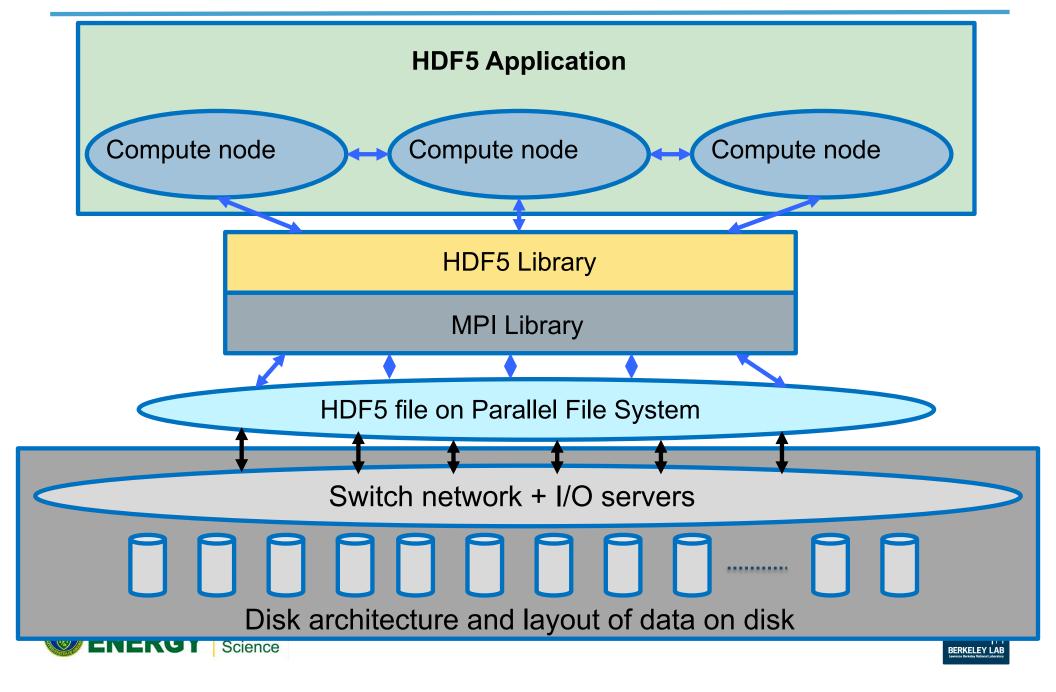
Standard HDF5 "Skeleton"

create (open) File **H5Fcreate (H5Fopen) H5Screate_simple/H5Screate_create_dataSpace H5Dcreate (H5Dopen)** create (open) Dataset H5Dread, H5Dwrite access Dataset **H5Dclose** close Dataset **H5Sclose** close dataSpace **H5Fclose** close File





PHDF5 Implementation Layers



Example of a PHDF5 C Program

A parallel HDF5 program has a few extra calls

```
MPI_Init(&argc, &argv);
fapl_id = H5Pcreate(H5P_FILE_ACCESS);
H5Pset_fapl_mpio(fapl_id, comm, info);
file_id = H5Fcreate(FNAME, ..., fapl_id);
space_id = H5Screate_simple(...);
dset_id = H5Dcreate(file_id, DNAME, H5T_NATIVE_INT,
                     space_id, ...);
xf_id = H5Pcreate(H5P_DATASET_XFER);
H5Pset_dxpl_mpio(xf_id, H5FD_MPIO_COLLECTIVE);
status = H5Dwrite(dset_id, H5T_NATIVE_INT, ..., xf_id...);
MPI_Finalize();
```





PHDF5 Etiquette

- PHDF5 opens a shared file with an MPI communicator
- Returns a file handle
- All future access to the file via that file handle
- All processes must participate in collective PHDF5 APIs
- Different files can be opened via different communicators
- <u>All HDF5 APIs that modify structural metadata are collective!</u>
 (file ops., object structure and life-cycle)

https://www.hdfgroup.org/HDF5/doc/RM/CollectiveCalls.html





Parallel HDF5 tutorial examples

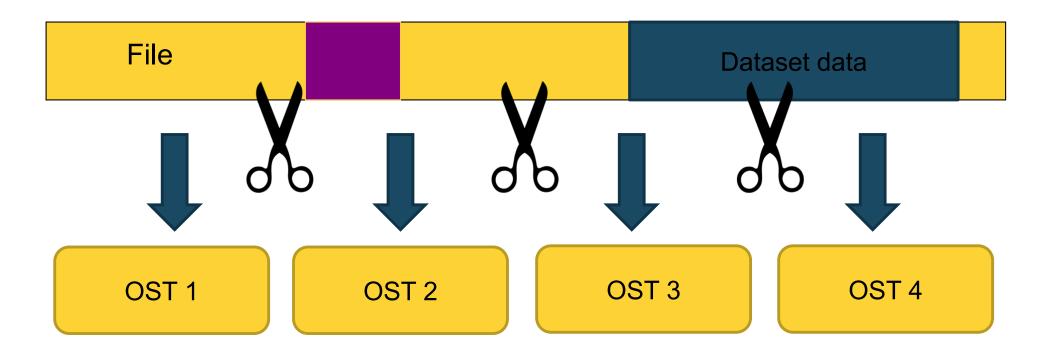
For simple examples how to write different data patterns see

http://www.hdfgroup.org/HDF5/Tutor/parallel.html





In a Parallel File System



The file is striped over multiple "disks" (e.g. Lustre OSTs) depending on the stripe size and stripe count with which the file was created.

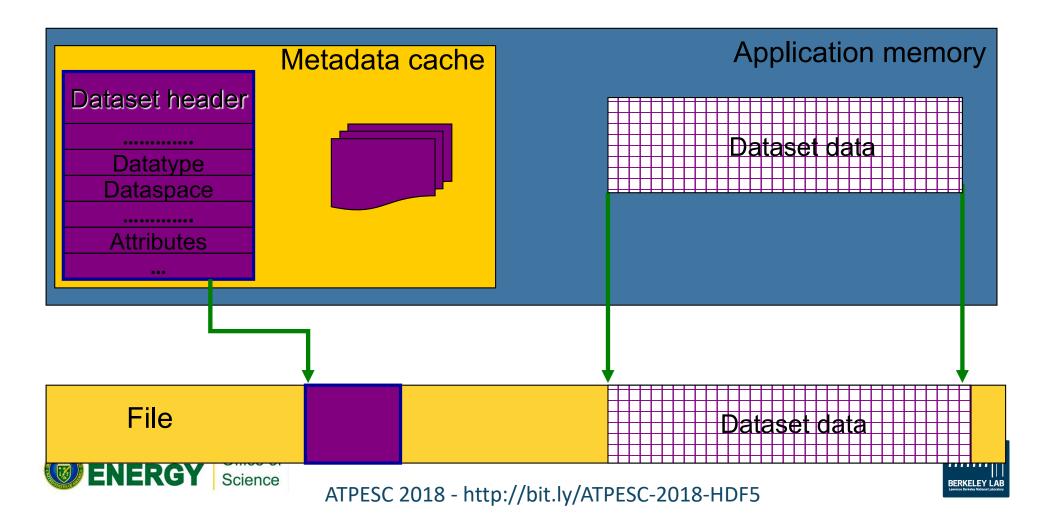
And it gets worse before it gets better...





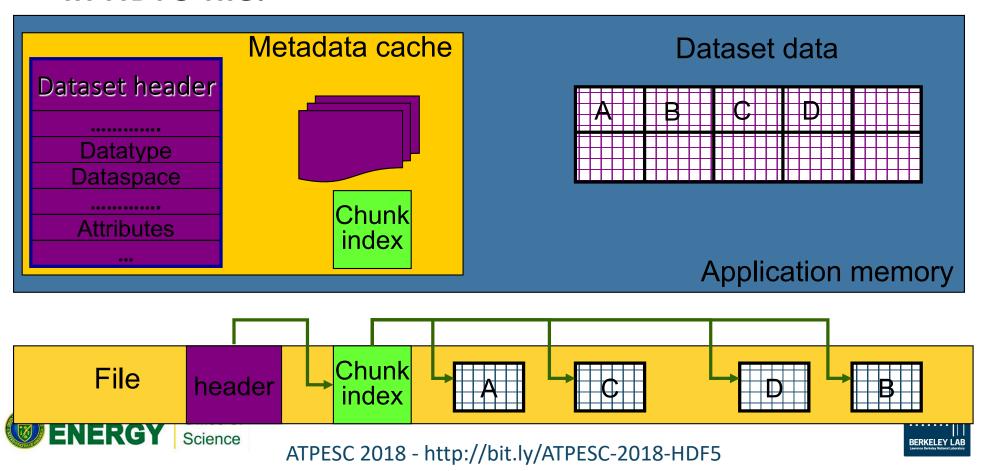
Contiguous Storage

- Metadata header separate from dataset data
- Data stored in one contiguous block in HDF5 file

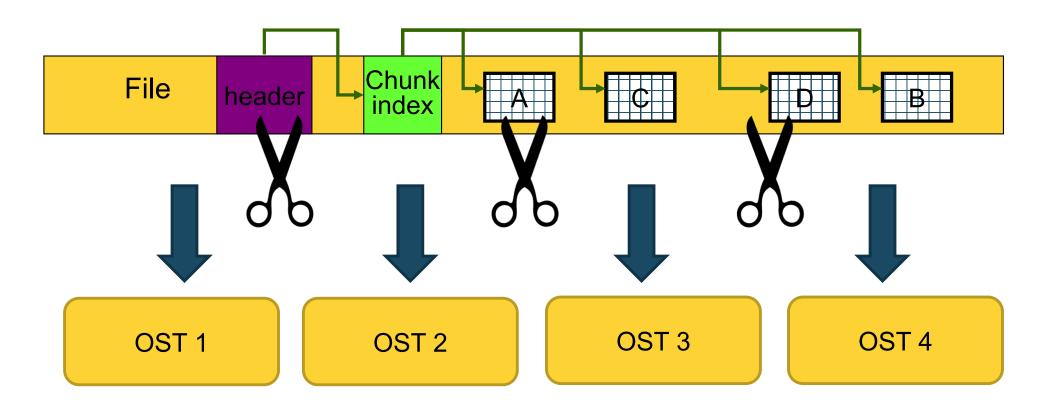


Chunked Storage

- Dataset data is divided into equally sized blocks (chunks).
- Each chunk is stored separately as a contiguous block in HDF5 file.



In a Parallel File System



The file is striped over multiple OSTs depending on the stripe size and stripe count with which the file was created.





Collective vs. Independent I/O

- Collective I/O attempts to combine multiple smaller independent I/O ops into fewer larger ops.
 - Neither mode is preferable a priori
- MPI definition of collective calls:
 - All processes of the communicator must participate in calls in the same order:

```
Process 1 Process 2 call A(); \rightarrow call B(); call A(); \rightarrow call B(); **right** call A(); \rightarrow call B(); call A(); **wrong**
```

- Independent calls are not collective [©]
- Collective calls are not necessarily synchronous, nor must they require communication
 - It could be that only internal state for the communicator changes



Data and Metadata I/O

Data

- Problem-sized
- I/O can be independent or collective
- Improvement targets:
 - Avoid unnecessary I/O
 - I/O frequency
 - Layout on disk
 - Different I/O strategies for chunked layout
 - Aggregation and balancing
 - Alignment

Metadata

- Small
- Reads can be independent or collective
- All modifying I/O must be collective
- Improvement targets:
 - Metadata design
 - Use the latest library version, if possible
 - Metadata cache
 - In desperate cases, take control of evictions





Don't Forget: It's a Multi-layer Problem

Application

HDF5

(Disable truncate in H5Fclose)

MPI-IO

(Number of collective buffer nodes, Collective buffer size)

Lustre Parallel File System

(Stripe factor and Stripe size)

Storage Hardware





Tools

DIAGNOSTICS AND INSTRUMENTATION



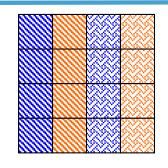


A Textbook Example

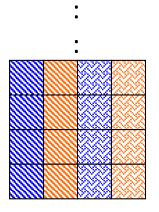
User reported:

 Independent data transfer mode is much slower than the collective data transfer mode

 Data array is tall and thin: 230,000 rows by 4 columns



230,000 rows







Symptoms

Writing to one dataset

- 4 MPI processes → 4 columns
- Datatype is 8-byte floats (doubles)
- 4 processes x 1000 rows x 8 bytes = 32,000 bytes
- % mpirun -np 4 ./a.out 1000
 - > Execution time: 1.783798 s.
- % mpirun -np 4 ./a.out 2000
 - > Execution time: 3.838858 s. (linear scaling)



2 sec. extra for 1000 more rows = 32,000 bytes.
 16KB/sec → Way too slow!!!





"Poor Man's Debugging"

- Build a version of PHDF5 with
- ./configure --enable-debug --enable-parallel ...
- This allows the tracing of MPIO I/O calls in the HDF5 library such as MPI_File_read_xx and MPI_File_write_xx
- Don't forget to % setenv H5FD_mpio_Debug "rw"
- You'll get something like this...





Independent and Contiguous

```
% setenv H5FD mpio Debug 'rw'
% mpirun -np 4 ./a.out 1000
                                 # Indep.; contiguous.
in H5FD_mpio_write mpi_off=0
                                size_i=96
                  mpi_off=0
                                size_i=96
in H5FD_mpio_write
                  mpi_off=0
in H5FD_mpio_write
                                size_i=96
                  mpi_off=0
                                size_i=96
in H5FD_mpio_write
                   mpi_off=2056
in H5FD_mpio_write
                                size_i=8
in H5FD_mpio_write
                   mpi_off=2048
                                size_i=8
                                size_i=8
in H5FD_mpio_write
                   mpi_off=2072
in H5FD_mpio_write
                   mpi_off=2064
                                size_i=8
                                size_i=8
in H5FD_mpio_write
                  mpi_off=2088
in H5FD_mpio_write
                   mpi_off=2080
                                size_i=8
```

A total of 4000 of these 8 bytes writes == 32,000 bytes.



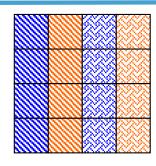


Plenty of Independent and Small Calls

Diagnosis:

 Each process writes one element of one row, skips to next row, writes one element, and so on.

Each process issues
 230,000 writes of 8
 bytes each.



: : 230,000 rows





Chunked by Column

```
% setenv H5FD mpio Debug 'rw'
% mpirun -np 4 ./a.out 1000
                                  # Indep., Chunked by column.
in H5FD_mpio_write mpi_off=0
                                size_i=96
in H5FD_mpio_write
                   mpi off=0
                                size i=96
                                               Metadata
                                size i=96
                   mpi off=0
in H5FD_mpio_write
                   in H5FD_mpio_write
in H5FD_mpio_write
                   mpi off=3688
                                    size_i=8000
                                    size_i=8000
in H5FD_mpio_write
                   mpi_off=11688
                                                  Dataset elements
in H5FD_mpio_write
                   mpi_off=27688
                                    size_i=8000
in H5FD_mpio_write
                   mpi_off=19688
                                    size_i=8000
in H5FD_mpio_write
                   mpi_off=96
                                    size_i=40
in H5FD_mpio_write
                   mpi off=136
                                    size i=544
                                                Metadata
in H5FD_mpio_write
                   mpi off=680
                                    size_i=120
                                    size_i=272
in H5FD_mpio_write
                   mpi_off=800
```

Execution time: 0.011599 s.



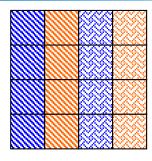


Use Collective Mode or Chunked Storage

Remedy:

 Collective I/O will combine many small independent calls into few but bigger calls

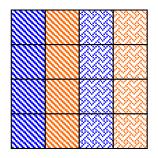
 Chunks of columns speeds up too



: : : 230,000 rows

:

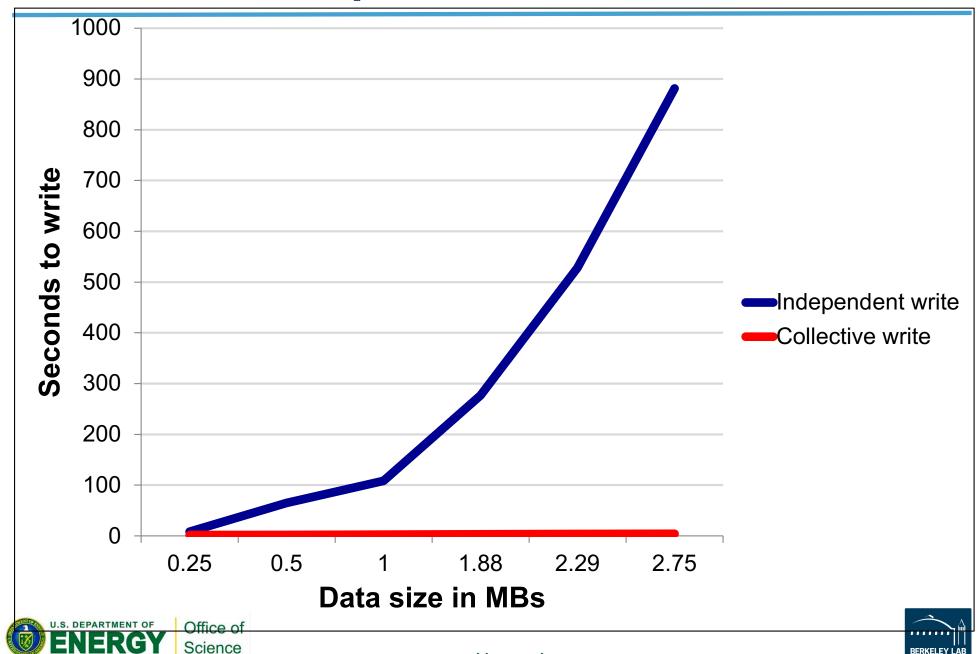
•







Collective vs. independent write



Back Into the Real World...

Two kinds of tools:

- I/O benchmarks for measuring a system's I/O capabilities
- I/O profilers for characterizing applications' I/O behavior

Two examples:

- h5perf (in the HDF5 source code distro)
- <u>Darshan</u> (from Argonne National Laboratory)

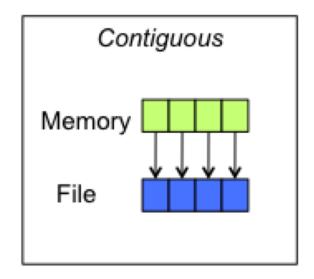
Profilers have to compromise between

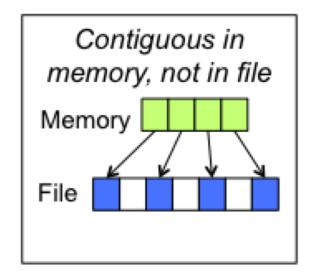
- A lot of detail → large trace files and overhead
- Aggregation → loss of detail, but low overhead

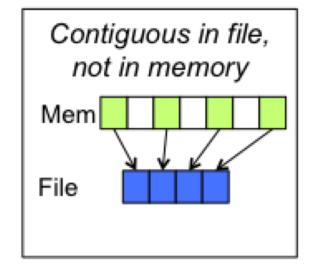


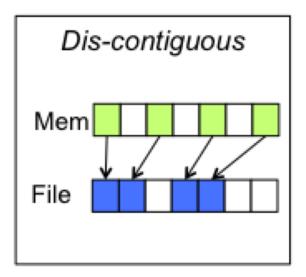


I/O Patterns













h5perf(_serial)

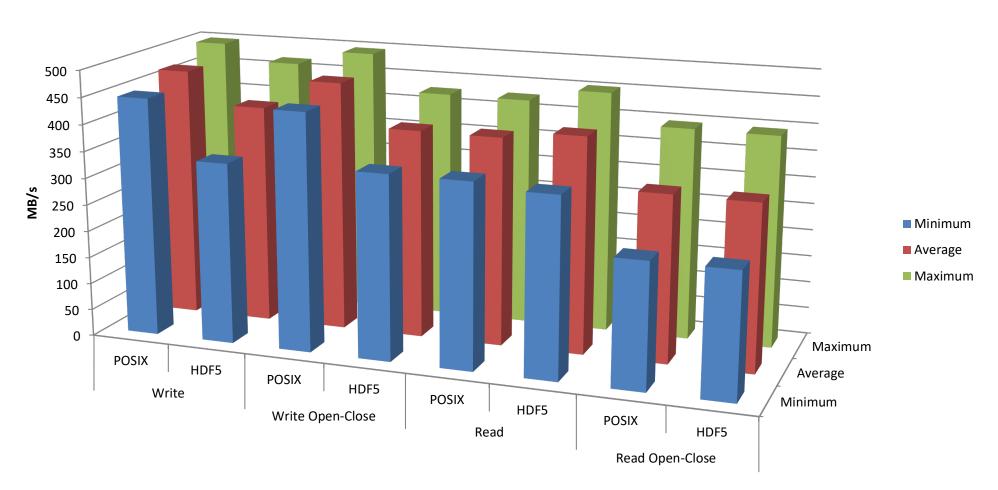
- Measures performance of a filesystem for different I/O patterns and APIs
- Three File I/O APIs for the price of one!
 - POSIX I/O (open/write/read/close...)
 - MPI-I/O (MPI_File_{open,write,read,close})
 - HDF5 (H5Fopen/H5Dwrite/H5Dread/H5Fclose)
- An indication of I/O speed ranges and HDF5 overheads
- Expectation management...





A Serial Run

h5perf_serial, 3 iterations, 1 GB dataset, 1 MB transfer buffer, HDF5 dataset contiguous storage, HDF5 SVN trunk, NCSA BW

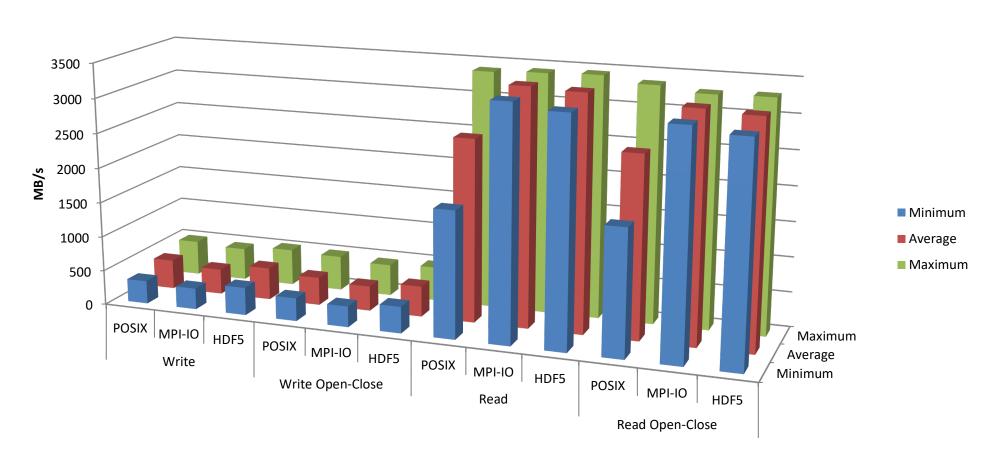






A Parallel Run

h5perf, 3 MPI processes, 3 iterations, 3 GB dataset (total), 1 GB per process, 1 GB transfer buffer, HDF5 dataset contiguous storage, HDF5 SVN trunk, NCSA BW







Darshan (ANL)

Design goals:

- Transparent integration with user environment
- Negligible impact on application performance

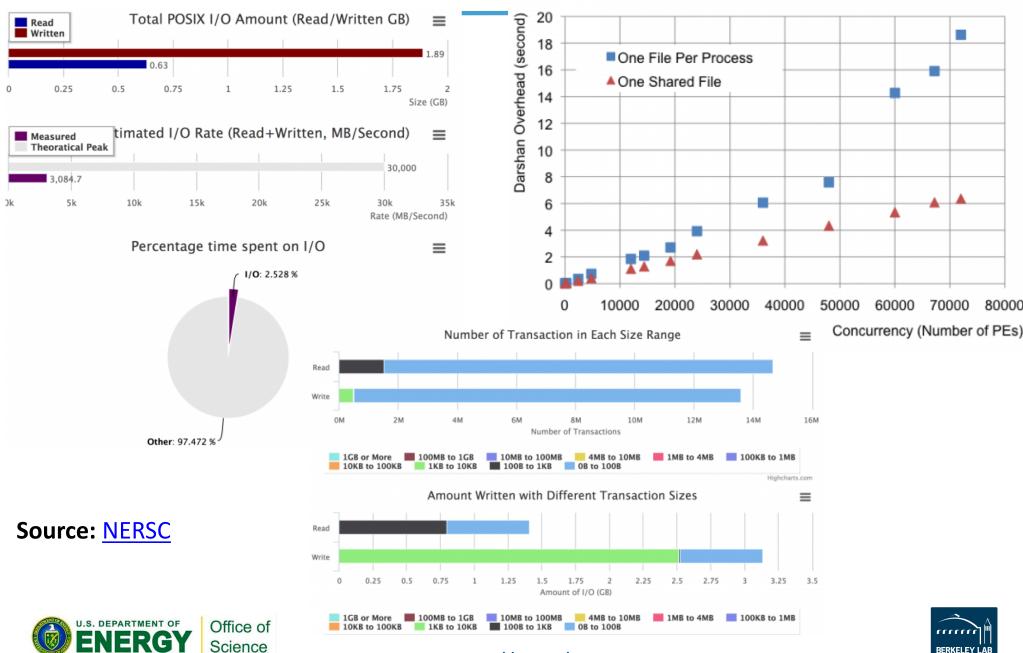
Provides aggregate figures for:

- Operation counts (POSIX, MPI-IO, HDF5, PnetCDF)
- Datatypes and hint usage
- Access patterns: alignments, sequentiality, access size
- Cumulative I/O time, intervals of I/O activity
- Does not provide I/O behavior over time
- Excellent starting point, maybe not your final stop





Darshan Sample Output



ATPESC 2018 - http://bit.ly/ATPESC-2018-HDF5

EXAMPLES





Standard Questions

- What I/O layers are involved and how much control do I have over them?
- Which ones do I tackle in which order?
 - Are there any low-hanging fruit?
- What's my baseline (for each layer) and what are my metrics?
- Which tool(s) will give me the information I need?
- When do I stop?
- New information → New answers (maybe): Need to keep an open mind!



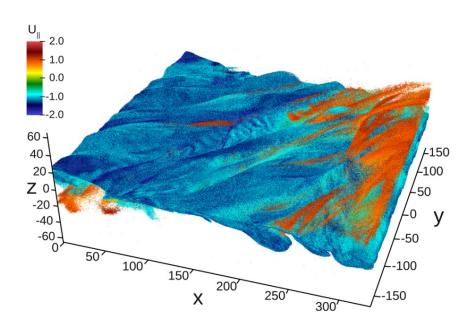


Reference:

<u>Trillion Particles, 120,000 cores, and 350 TBs:</u>
<u>Lessons Learned from a Hero I/O Run on Hopper,</u>
By Suren Byna (LBNL) et al., 2015.

Examples

VPIC







Layers

Application

HDF5

(Disable truncate in H5Fclose)

MPI-IO

(Number of collective buffer nodes, Collective buffer size)

Lustre Parallel File System

(Stripe factor and Stripe size)

Storage Hardware





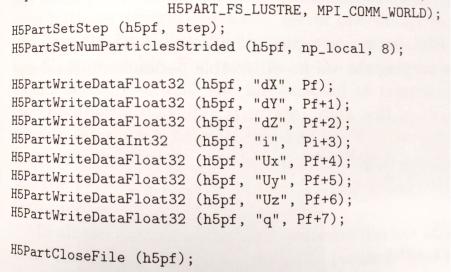
"Application I/O Structure"

- Total control over all layers
- Challenge: large output files
- Metric: write speed (throughput)
- Computationally intensive

 Need an I/O kernel
- H5Part multiple dataset writes

- "Game plan":
 - MPI-IO / Lustre tuning
 - Low hanging fruit (relatively)
 - Pair MPI aggregators with Lustre OSTs
 - Match MPI-IO buffer sizes and Lustre stripe size
 - Worry about HDF5 (H5Part)





h5pf = H5PartOpenFileParallel (fname, H5PART_WRITE



I/O Aggregation

Application
MPI Domains
20,000
6 threads per domain

MPI Aggregators

Aggregator

Aggregator

OST

OST

OST

OST



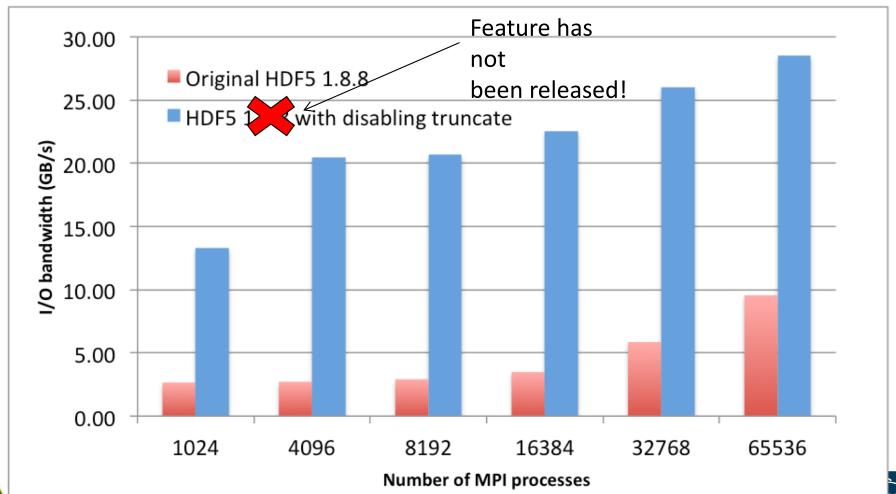




Closing HDF5 File ...

Q: How long does it take to close/flush an HDF5 file?

A: A lot longer than you might expect!

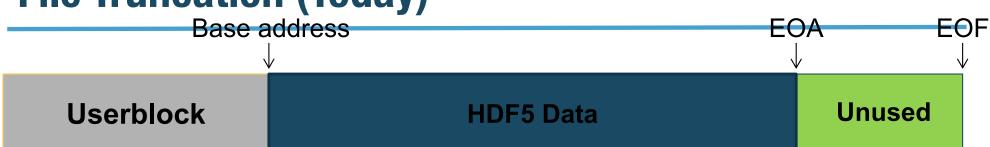




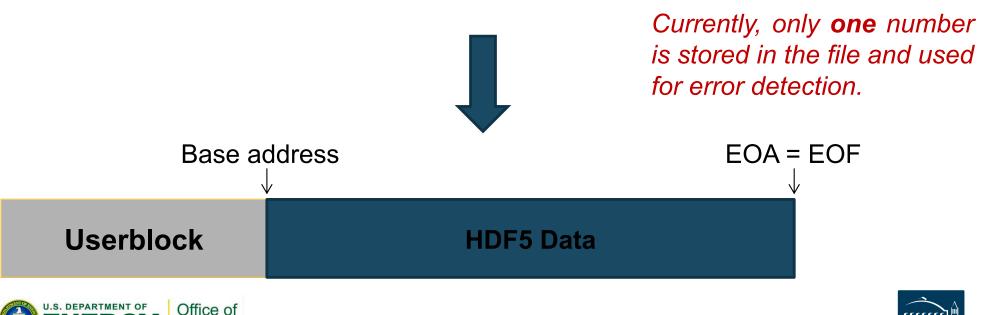
Science

File Truncation (Today)

Science



A call to H5Fflush or H5Fclose triggers a call to ftruncate (serial) or MPI_File_set_size (parallel), which can be fairly expensive.





File Truncation (Tomorrow)



A call to H5Ff1ush or H5Fc1ose triggers both values (EOA, EOF) to be saved in the file and **no** truncation takes place, IF the file was created with the "avoid truncation" property set.



Caveat: Incompatible with older versions of the library. Requires HDF5 library version 1.12 or later.





Multi-Dataset I/O - Motivation

- HDF5 accesses elements in one dataset at a time
- Many HPC applications access data in multiple datasets in every time step
- Frequent small-size dataset access → Big Trouble (≠Big Data)
- Parallel file systems tend not to like that.
- Idea: Let users to do more I/O per HDF5 call!
- Two New API routines:
 - H5Dread_multi()
 - H5Dwrite_multi()

Not a new idea: PnetCDF has supported this for some time...

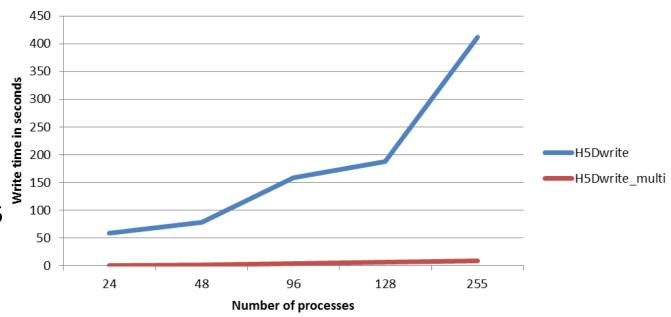




Sample Results

The plot shows the performance difference between using a single H5Dwrite() multiple times and using H5Dwrite_multi() once on 30 chunked datasets.

(On Hopper @ NERSC, a Cray XE-6 with Lustre file system)



Run	Code	Nodes	Cores	File Size	Stripe Size	Write Time	Total HDF Time	Throughput
Max Throughput	Multiple Dataset	320	5,120	5 TB	1 GB	91.08 s	167.78 s	56.21 GB/s
Code Comparison	Single Dataset	320	5,120	5 TB	128 MB	116.94 s	117.37 s	43.78 GB/s
Hero Run	Single Dataset	9,314	298,048	291 TB	1 GB	5,763.14 s	5,779.89 s	51.81 GB/s

TABLE I

COMPARISON OF VPIC-IO KERNEL PARAMETERS AND OBSERVED IO THROUGHPUT.





Reference:

<u>Parallel and Large-scale Simulation Enhancements to CGNS</u>, By Scot Breitenfeld, The HDF Group, 2015.

Examples

CGNS







CFD Standard

- CGNS = Computational Fluid Dynamics (CFD) General Notation System
- An effort to standardize CFD input and output data including:
 - Grid (both structured and unstructured), flow solution
 - Connectivity, boundary conditions, auxiliary information.

Two parts:

- A standard format for recording the data
- Software that reads, writes, and modifies data in that format.
- An American Institute of Aeronautics and Astronautics Recommended Practice





CGNS Storage Evolution

- CGNS data was originally stored in ADF ('Advanced Data Format')
- ADF lacks parallel I/O or data compression capabilities
- Doesn't have HDF5's support base and tools
- HDF5 superseded ADF as the official storage mechanism
- CGNS introduced parallel I/O APIs w/ parallel HDF5 in 2013
- Poor performance of the new parallel APIs in most circumstances
- In 2014, NASA provided funding for The HDF Group with the goal to improve the under-performing parallel capabilities of the CGNS library.





CGNS Performance Problems

Opening an existing file

- CGNS reads the entire HDF5 file structure, loading a lot of (HDF5) metadata
- Reads occur independently on ALL ranks competing for the same metadata
 - →"Read Storm"

Closing a CGNS file

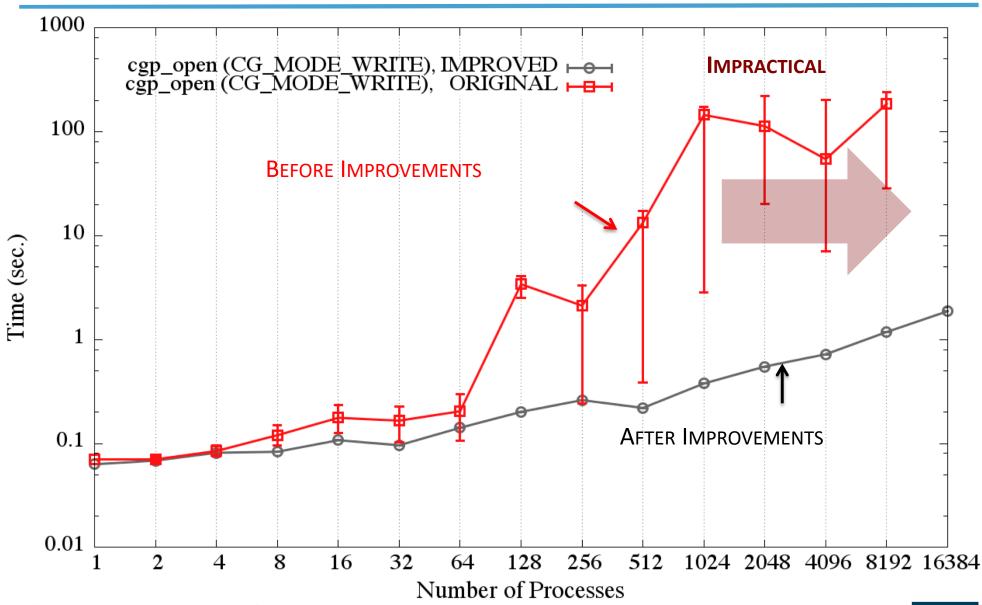
- Triggers HDF5 flush of a large amount of small metadata entries
- Implemented as iterative, independent writes, an unsuitable workload for parallel file systems







Opening CGNS File





Metadata Read Storm Problem (I)

All metadata "write" operations are required to be collective:

```
if(0 == rank)
    H5Dcreate("dataset1");
else if(1 == rank)
    H5Dcreate("dataset2");
```

```
/* All ranks have to call */
H5Dcreate("dataset1");
H5Dcreate("dataset2");
```

 Metadata read operations are not required to be collective

```
if(0 == rank)
    H5Dopen("dataset1");
else if(1 == rank)
    H5Dopen("dataset2");
```

```
/* All ranks have to call */
H5Dopen("dataset1");
H5Dopen("dataset2");
```





Metadata Read Storm Problem (II)

- Metadata read operations are treated by the library as independent read operations.
- Consider a very large MPI job size where all processes want to open a dataset that already exists in the file.
- All processes
 - Call H5Dopen("/G1/G2/D1");
 - Read the same metadata to get to the dataset and the metadata of the dataset itself
 - IF metadata not in cache, THEN read it from disk.
 - Might issue read requests to the file system for the same small metadata.
- → READ STORM





Avoiding a Read Storm

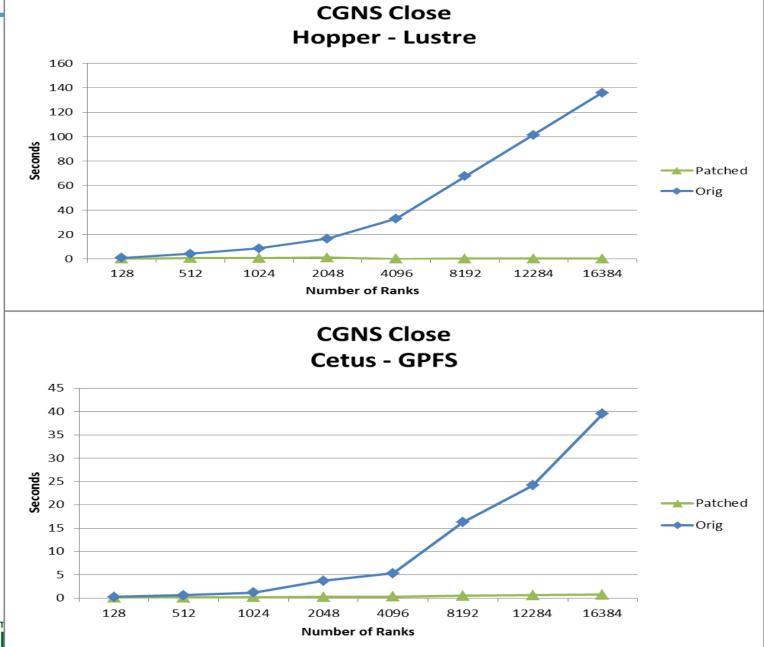
- Hint that metadata access is done collectively
- A property on an access property list
- If set on the file access property list, then all metadata read operations will be required to be collective
- Can be set on individual object property list
- If set, MPI rank 0 will issue the read for a metadata entry to the file system and broadcast to all other ranks







Closing a CGNS File ...







Write Metadata Collectively!

- Symptoms: Many users reported that H5Fclose() is very slow and doesn't scale well on parallel file systems.
- Diagnosis: HDF5 metadata cache issues very small accesses (one write per entry). We know that parallel file systems don't do well with small I/O accesses.
- Solution: Gather up all the entries of an epoch, create an MPI derived datatype, and issue a single collective MPI write.





HDF5 Roadmap

- Concurrency
 - Single-Writer / Multiple-Reader (SWMR)
 - Asynchronous I/O
 - Internal threading
- Virtual Object Layer
- Virtual Datasets
- Query & Indexing
- Native HDF5 client/server

- Performance
 - Scalable chunk indices
 - Metadata aggregation and Page buffering
 - Variable-length records
- Fault tolerance
- Parallel I/O
- I/O Autotuning





Questions, Comments, Feedback?

Thank You!





ExaHDF5 - Features

- Virtual Object Layer (VOL) integration into HDF5
- Caching and prefetching Data Elevator VOL
- Topology-aware I/O
- Asynchronous I/O
- Independent metadata updates
- Workflow supporting features SWMR
- Querying HDF5 data and metadata
- Interoperability with other file formats
 - PnetCDF/netCDF, ADIOS
- Maintenance and release support
- ECP engagement w/ AD, ST, and HT
 - Consulting and performance tuning for applications





ExaHDF5 – Development timeline

